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# Satoku Matrix Oreganography

Release 0.0

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## Abstract

Oreganography is the art of encrypting information in weedy ways for herbal recreation.

## Contents

List of Figures	1
List of Tables	2
1 Satoku Matrix	2
2 Mapping a CNF Formula to a Satoku Matrix	4
3 Mapping a Satoku Matrix to a CNF Formula	5
4 Initial Information Encoding	6
5 Advanced Information Hiding	9
6 Satoku Matrix Mapped to CNF Formula	16
7 CNF Formula Without At-Least-1 Clauses	17

## List of Figures

1.1 Satoku Matrix Index Scheme Example	2
1.2 Merge State Rows $S$ Into State Row $s_{i_j}$	3
1.3 Requirement Update Algorithm	3
1.4 Consolidated Satoku Matrix Index Scheme Example	3
2.1 Mapping a CNF Formula to a Satoku Matrix	4
3.1 Mapping a Satoku Matrix to a CNF Formula	5
4.1 Text Encoded in a Satoku Matrix	6
4.2 Reduced CNF Formula of Mapped Satoku Matrix	6
4.3 Reduced CNF Formula Mapped to Unconsolidated Satoku Matrix	7
4.4 Consolidated Satoku Matrix for Reduced CNF Formula	8
5.1 Reduced CNF Formula Without At-Least-1 Clauses	9
5.2 Consolidated Satoku Matrix for Reduced CNF Formula Without At-Least-1 Clauses	9
5.3 Unconsolidated Satoku Matrix of Reduced CNF With Mapped Variables	10
5.4 Consolidated Satoku Matrix of Reduced CNF With Mapped Variables	11
5.5 Mapped Variable Quadrant of Full CNF Formula Without At-Least-1 Clauses	12
5.6 At-Least-1 Clauses Manually Added to Mapped Variable Quadrant	13
5.7 Requirement Update Restores Original Message	14

5.8	Removing Mapped Variables Leaves Exact Copy of Original Message	15
6.1	Satoku Matrix Mapped to CNF Formula	16
7.1	CNF formula without at-least-1 clauses	17

## List of Tables

1.1	Satoku Matrix Index Scheme	2
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## 1 Satoku Matrix

A satoku matrix is an inverted adjacency matrix preserving clause boundary information for implementing a requirement update algorithm (see figure 1.3).

A satoku matrix  $\mathbb{S}$  is formally defined as a sequence of *cell-matrix rows*  $c_i$ ,  $0 \leq i < |\mathbb{S}|$ , where a *cell-matrix row*  $c_i$  consists of *cells*  $c_{i_g}$ ,  $0 \leq g < |c_i|$ . A *cell*  $c_{i_g}$  consists of *cell rows*  $r_{i_{jg}}$ ,  $0 \leq j < |c_{i_g}|$ , where a *cell row*  $r_{i_{jg}}$  consists of *states*  $s_{i_{jgh}}$ ,  $0 \leq h < |r_{i_{jg}}|$ . A *state*  $s_{i_{jgh}}$  represents a *conflict relationship* CFR with the *state*  $s_{gh_{i_j}}$ , where a *conflict relationship* is either *possible* = 1 or *impossible* = 0. A *state row*  $s_{i_j}$ ,  $0 \leq j < |c_i|$  is the sequence of *cell rows*  $r_{i_{jg}}$ . The index scheme is summarized in table 1.1.

table 1.1: Satoku Matrix Index Scheme

indexed state entity	description
$C$ <i>cell-matrix-row</i>	row $c_i$ of cells (cell-matrix-row)
$C$ <i>cell-matrix-row</i> <i>cell-matrix-column</i>	single cell $c_{i_g}$
$r$ <i>cell-matrix-row</i> <i>cell-row</i> <i>cell-matrix-column</i>	cell row $r_{i_{jg}}$ containing all CFR states between an atomic state $s_{i_{jj}}$ and an atomic cell $c_{i_g}$
$S$ <i>cell-matrix-row</i> <i>cell-row</i>	state row $s_{i_j}$ of all cell rows $r_{i_{jg}}$ containing all singular states for an atomic state $s_{i_{jj}}$
$S$ <i>cell-matrix-row</i> <i>cell-row</i> <i>cell-matrix-column</i> <i>cell-column</i>	singular state $s_{i_{jgh}}$

For better readability, if a *cell row*  $r_{i_{jg}}$  contains more than one 1-state, all 1-states in  $r_{i_{jg}}$  are represented by a dash (-) (see figure 1.1).

P	---	----	---	--	
$s_{00}$	1 0 0	----	---	--	$c_0$
$s_{01}$	0 1 0	0----	---	--	
$s_{02}$	0 0 1	-----	-0-	0 1	
$s_{10}$	-0-	1 0 0 0	0--	--	$c_{12}$
$s_{11}$	---	0 1 0 0	-0-	0 1	
$s_{12}$	---	0 0 1 0	---	1 0	
$s_{13}$	---	0 0 0 1	--0	--	
$s_{20}$	---	0----	1 0 0	--	$s_{20}$
$s_{21}$	--0	-0--	0 1 0	--	$s_{210}$
$s_{22}$	---	---0	0 0 1	--	$s_{2213}$
$s_{30}$	--0	-0--	---	1 0	$r_{312}$
$s_{31}$	---	--0-	---	0 1	

figure 1.1: Satoku Matrix Index Scheme Example

Merging a sequence of *state rows*  $S$  into a *state row*  $s_{ij}$ , denoted as  $\text{Mrg}(s_{ij}, S)$ , is defined by the algorithm in figure 1.2. It returns the number of  $1 \rightarrow 0$  transitions performed.

---

```

def Mrg( $s_{ij}, S$ ):
    transitions  $\leftarrow$  0
    for each state row  $s_{gh} \in S$ :
        for each state  $s_{i_{jef}}, 0 \leq e < |S|, 0 \leq f < |r_{i_{je}}|$ :
            if  $s_{i_{jef}} \neq 0$  and  $s_{gh_{ef}} = 0$ :
                 $s'_{i_{jef}} = 0$ , transitions  $+=$  1
            if  $s_{ef_{ij}} \neq 0$ :
                 $s'_{ef_{ij}} = 0$ , transitions  $+=$  1
    return transitions

```

---

figure 1.2: Merge State Rows  $S$  Into State Row  $s_{ij}$

The requirement update algorithm in figure 1.3 distributes conflict relationships into all *cell rows* which have a single 1-state. After applying the requirement update algorithm, the satoku matrix  $\mathbb{S}$  is called *consolidated*.

---

```

transitions  $\leftarrow$  1
while transitions  $>$  0:
    transitions  $\leftarrow$  0
    for each state row  $s_{ij}$ :
        for each cell row  $r_{ijg}, i \neq g$ :
            if there is only a single 1-state  $s_{ij_{gh}} = 1, \sum_{f=0}^{|r_{ijg}|-1} s_{ij_{gf}} = 1$ :
                transitions  $+=$   $\text{Mrg}(s_{ij}, s_{gh})$  // merge state row  $s_{gh}$  into state row  $s_{ij}$ 

```

---

figure 1.3: Requirement Update Algorithm

The consolidated version of the Satoku Matrix Index Scheme Example from figure 1.1 is shown in figure 1.4. Notably  $s_{0_23_1}$  triggered a merge of  $s_{3_1}$  into  $s_{0_2}$ . This led to state  $s_{3_{1_2}}$  causing  $1 \rightarrow 0$  transitions of state  $s_{0_{2_{1_2}}}$  and state  $s_{1_{2_{0_2}}}$ .

P	---	----	---	--
$s_{0_0}$	1 0 0	-----	----	--
$s_{0_1}$	0 1 0	0----	----	--
$s_{0_2}$	0 0 1	--0--	-0-	0 1
$s_{1_0}$	-0-	1 0 0 0	0--	--
$s_{1_1}$	---	0 1 0 0	-0-	0 1
$s_{1_2}$	--0	0 0 1 0	---	1 0
$s_{1_3}$	---	0 0 0 1	--0	--
$s_{2_0}$	---	0----	1 0 0	--
$s_{2_1}$	--0	-0--	0 1 0	--
$s_{2_2}$	---	---0	0 0 1	--
$s_{3_0}$	--0	-0--	---	1 0
$s_{3_1}$	---	--0-	---	0 1

figure 1.4: Consolidated Satoku Matrix Index Scheme Example

## 2 Mapping a CNF Formula to a Satoku Matrix

A CNF formula  $F$  is a conjunction of  $m$  disjunctive clauses  $C_i$  each containing  $k_i$  literals  $l_j$ , where a literal  $l_j$  is a negated or unnegated boolean variable:

$$F = \bigwedge_{i=0}^{m-1} C_i, \quad m = |F|, \quad C_i = \bigvee_{j=0}^{k_i-1} l_j, \quad k_i = |C_i|, \quad m, k_i \in \mathbb{N}_0.$$

Mapping a CNF formula  $F$  to a satoku matrix  $\mathbb{S}$  with the algorithm in figure 2.1 results in an *unconsolidated* satoku matrix  $\mathbb{S}$ .

---

```

for each clause  $C_i$  of CNF formula  $F$ ,  $0 \leq i < |F|$ :
  extend each state row  $s_{e_f}$ , by  $|C_i|$  1-states,  $0 \leq e < i$ ,  $0 \leq f < |c_{e_e}|$ 
  add  $|C_i|$  state rows with  $\sum_{n=0}^i |C_n|$  1-states each
  // process at-most-1 constraints
  for each cell row  $r_{i_{j_i}}$ ,  $0 \leq j < |C_i| - 1$ :
    for each state  $s_{i_{j_{i_h}}}$ ,  $j < h < |C_i|$ :
       $s_{i_{j_{i_h}}} \leftarrow 0$ 
       $s_{i_{h_{i_j}}} \leftarrow 0$ 
  // process conflict relationships
  for each state row  $s_{i_j}$ ,  $0 \leq i < |F|$ ,  $0 \leq j < |C_i|$ :
    for each cell row  $r_{i_{j_g}}$ ,  $i < g < |F|$ :
      for each state  $s_{i_{j_{g_h}}}$ ,  $0 \leq h < |C_g|$ :
        if  $C_{i_j} \wedge C_{g_h} = 0$ :
           $s_{i_{j_{g_h}}} \leftarrow 0$ 
           $s_{g_{h_{i_j}}} \leftarrow 0$ 

```

---

figure 2.1: Mapping a CNF Formula to a Satoku Matrix

### 3 Mapping a Satoku Matrix to a CNF Formula

See figure 3.1 for an algorithm to map a satoku matrix to a CNF formula.

---

```
start with empty CNF formula  $F$  (a conjunctive clause)
for each cell  $c_{i_i}$ ,  $0 \leq i < |\mathbb{S}|$ :
    add an empty disjunctive clause  $C_i$  to  $F$ 
    for each cell row  $r_{i_i}$ ,  $0 \leq j < |c_{i_i}|$ :
        add an unnegated variable  $v_{i_j}$  to clause  $C_i$ 
for each state row  $s_{i_j}$ ,  $0 \leq i < |\mathbb{S}|$ ,  $0 \leq j < |c_{i_i}|$ :
    with cell row  $r_{i_i}$ :
        for each state  $s_{i_j i_f}$ ,  $j < f < |r_{i_i}|$ :
            if  $s_{i_j i_f} = 0$ :
                to express the conclusion  $v_{i_j} \rightarrow \neg v_{i_f}$ ,
                add the disjunctive clause  $\neg v_{i_j} \vee \neg v_{i_f}$  to  $|F|$ 
for each cell row  $r_{i_j e}$ ,  $i < e < |\mathbb{S}|$ :
    for each state  $s_{i_j e f}$ ,  $0 \leq f < |r_{i_j e}|$ :
        if  $s_{i_j e f} = 0$ :
            add the disjunctive clause  $\neg v_{i_j} \vee \neg v_{e_f}$  to  $|F|$ 
```

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figure 3.1: Mapping a Satoku Matrix to a CNF Formula

## 4 Initial Information Encoding

Arbitrary pixel blocks like “HAPPYEASTER!<smile><egg><egg>” can be encoded in the upper right half of a satoku matrix in a simple manner by flipping 1-states to 0 (see figure 4.1).

P	-----	-----	-----	-----	-----	-----
s <sub>00</sub>	1 0 0 0 0 0	- 0 - - 0 -	- - 0 0 - -	- 0 0 0 - -	- 0 0 0 - -	- 0 - - - 0
s <sub>01</sub>	0 1 0 0 0 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - 0 -	- 0 - - 0 -	- - 0 - 0 -
s <sub>02</sub>	0 0 1 0 0 0	- 0 0 0 0 -	- 0 0 0 0 -	- 0 0 0 - -	- 0 0 0 - -	- - - 0 - -
s <sub>03</sub>	0 0 0 1 0 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - - -	- 0 - - - -	- - - 0 - -
s <sub>04</sub>	0 0 0 0 1 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - - -	- 0 - - - -	- - - 0 - -
s <sub>05</sub>	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -	- - - - - -	- - - - - -
s <sub>10</sub>	- - - - - -	1 0 0 0 0 0	- - 0 0 - -	- - 0 0 - -	- 0 0 0 0 -	1 0 0 0 0 0
s <sub>11</sub>	0 0 0 0 0 1	0 1 0 0 0 0	- 0 - - - -	- 0 - - 0 -	- 0 - - - -	- - - 0 - -
s <sub>12</sub>	- - 0 - - -	0 0 1 0 0 0	- 0 0 - - -	- 0 0 0 0 -	- 0 0 0 0 -	- - - 0 - -
s <sub>13</sub>	- 0 - - - -	0 0 0 1 0 0	- 0 - - - -	- 0 - - 0 -	- - - - 0 -	- - - 0 - -
s <sub>14</sub>	0 0 0 0 0 1	0 0 0 0 1 0	- - 0 0 - -	- 0 - - 0 -	- 0 0 0 0 -	- - - 0 - -
s <sub>15</sub>	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -	- - - - - -
s <sub>20</sub>	- - - - - -	- - - - - -	1 0 0 0 0 0	- - 0 0 - -	- 0 0 0 - -	- - - 0 - -
s <sub>21</sub>	- 0 0 0 0 -	- 0 0 0 - -	0 1 0 0 0 0	- 0 - - - -	- 0 - - 0 -	- - - 0 - -
s <sub>22</sub>	0 - 0 - - -	0 - 0 - 0 -	0 0 1 0 0 0	- 0 0 - - -	- 0 0 0 - -	- - - 0 - -
s <sub>23</sub>	0 - 0 - - -	0 - - - 0 -	0 0 0 1 0 0	- 0 - - - -	- 0 - 0 - -	- - - - - -
s <sub>24</sub>	- 0 0 0 0 -	- - - - - -	0 0 0 0 1 0	- - 0 0 - -	- 0 - - 0 -	- - - 0 - -
s <sub>25</sub>	- - - - - -	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -
s <sub>30</sub>	- - - - - -	- - - - - -	- - - - - -	1 0 0 0 0 0	- - - - - -	- - - - - -
s <sub>31</sub>	0 0 0 0 0 1	- 0 0 0 0 -	- 0 0 0 - -	0 1 0 0 0 0	- 0 - - 0 -	- - 0 0 - -
s <sub>32</sub>	0 - 0 - - -	0 - 0 - - -	0 - 0 - 0 -	0 0 1 0 0 0	- - - - - -	- 0 - - 0 -
s <sub>33</sub>	0 - 0 - - -	0 - 0 - - -	0 - - - 0 -	0 0 0 1 0 0	0 - - - - 0	- 0 - - 0 -
s <sub>34</sub>	- 0 - - - -	- 0 0 0 0 -	- - - - - -	0 0 0 0 1 0	- 0 0 0 0 -	- - 0 0 - -
s <sub>35</sub>	- - - - - -	- - - - - -	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -
s <sub>40</sub>	- - - - - -	- - - - - -	- - - - - -	- - - 0 - -	1 0 0 0 0 0	- - - - - -
s <sub>41</sub>	0 0 0 0 0 1	0 0 0 - 0 -	0 0 0 0 0 1	- 0 - - 0 -	0 1 0 0 0 0	- - 0 0 - -
s <sub>42</sub>	0 - 0 - - -	0 - 0 - 0 -	0 - 0 - - -	- - - - 0 -	0 0 1 0 0 0	- 0 - - 0 -
s <sub>43</sub>	0 - 0 - - -	0 - 0 - 0 -	0 - 0 0 - -	- - - - 0 -	0 0 0 1 0 0	- 0 - - 0 -
s <sub>44</sub>	- 0 - - - -	0 - 0 0 0 -	- 0 - - 0 -	- 0 - - 0 -	0 0 0 0 1 0	- - 0 0 - -
s <sub>45</sub>	- - - - - -	- - - - - -	- - - - - -	- - - 0 - -	0 0 0 0 0 1	- - - - - -
s <sub>50</sub>	- - - - - -	- - - - - -	- - - - - -	- - - - - -	- - - - - -	1 0 0 0 0 0
s <sub>51</sub>	0 - - - - -	0 - - - - -	- - - - - -	- - 0 0 - -	- - 0 0 - -	0 1 0 0 0 0
s <sub>52</sub>	- 0 - - - -	0 - - - - -	- - - - - -	- 0 - - 0 -	- 0 - - 0 -	0 0 1 0 0 0
s <sub>53</sub>	- - 0 0 0 -	0 0 0 0 0 1	0 0 0 - 0 -	- 0 - - 0 -	- 0 - - 0 -	0 0 0 1 0 0
s <sub>54</sub>	- 0 - - - -	0 - - - - -	- - - - - -	- - 0 0 - -	- - 0 0 - -	0 0 0 0 1 0
s <sub>55</sub>	0 - - - - -	0 - - - - -	- - - - - -	- - - - - -	- - - - - -	0 0 0 0 0 1

figure 4.1: Text Encoded in a Satoku Matrix

The satoku matrix in figure 4.1 can then be mapped to CNF formula (see figure 6.1). Since this formula is huge, only a reduced version is shown here in figure 4.2.

$$\begin{aligned}
 & ( v_{00} \vee v_{01} \vee v_{02} \vee v_{03} \vee v_{04} \vee v_{05} ) \wedge ( v_{10} \vee v_{11} \vee v_{12} \vee v_{13} \vee v_{14} \vee v_{15} ) \wedge ( \neg v_{00} \vee \neg v_{01} ) \wedge \\
 & ( \neg v_{00} \vee \neg v_{02} ) \wedge ( \neg v_{00} \vee \neg v_{03} ) \wedge ( \neg v_{00} \vee \neg v_{04} ) \wedge ( \neg v_{00} \vee \neg v_{05} ) \wedge ( \neg v_{00} \vee \neg v_{11} ) \wedge ( \neg v_{00} \vee \neg v_{14} ) \wedge ( \neg v_{01} \vee \neg v_{02} ) \wedge ( \neg v_{01} \vee \neg v_{03} ) \wedge \\
 & ( \neg v_{01} \vee \neg v_{04} ) \wedge ( \neg v_{01} \vee \neg v_{05} ) \wedge ( \neg v_{01} \vee \neg v_{11} ) \wedge ( \neg v_{01} \vee \neg v_{14} ) \wedge ( \neg v_{02} \vee \neg v_{03} ) \wedge ( \neg v_{02} \vee \neg v_{04} ) \wedge ( \neg v_{02} \vee \neg v_{05} ) \wedge ( \neg v_{02} \vee \neg v_{11} ) \wedge \\
 & ( \neg v_{02} \vee \neg v_{12} ) \wedge ( \neg v_{02} \vee \neg v_{13} ) \wedge ( \neg v_{02} \vee \neg v_{14} )
 \end{aligned}$$

figure 4.2: Reduced CNF Formula of Mapped Satoku Matrix





## 5 Advanced Information Hiding

Removing the at-least-1 clauses from the reduced CNF formula in figure 4.3 results in a set of 2-literal clauses (see figure 5.1). For the full formula in figure 6.1 without at-least-1 clauses see figure 7.1.

$$\begin{aligned}
 &(\neg v_{00} \vee \neg v_{01}) \wedge (\neg v_{00} \vee \neg v_{02}) \wedge (\neg v_{00} \vee \neg v_{03}) \wedge (\neg v_{00} \vee \neg v_{04}) \wedge (\neg v_{00} \vee \neg v_{05}) \wedge (\neg v_{00} \vee \neg v_{11}) \wedge (\neg v_{00} \vee \neg v_{14}) \wedge (\neg v_{01} \vee \neg v_{02}) \wedge \\
 &(\neg v_{01} \vee \neg v_{03}) \wedge (\neg v_{01} \vee \neg v_{04}) \wedge (\neg v_{01} \vee \neg v_{05}) \wedge (\neg v_{01} \vee \neg v_{11}) \wedge (\neg v_{01} \vee \neg v_{14}) \wedge (\neg v_{02} \vee \neg v_{03}) \wedge (\neg v_{02} \vee \neg v_{04}) \wedge (\neg v_{02} \vee \neg v_{05}) \wedge \\
 &(\neg v_{02} \vee \neg v_{11}) \wedge (\neg v_{02} \vee \neg v_{12}) \wedge (\neg v_{02} \vee \neg v_{13}) \wedge (\neg v_{02} \vee \neg v_{14})
 \end{aligned}$$

figure 5.1: Reduced CNF Formula Without At-Least-1 Clauses

Producing a consolidated satoku matrix from the formula in figure 5.1 lacks any information whatsoever (see figure 5.2).

P	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>00</sub>	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>01</sub>	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>10</sub>	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>11</sub>	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>20</sub>	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>21</sub>	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>30</sub>	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>31</sub>	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>40</sub>	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>41</sub>	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>50</sub>	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>51</sub>	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>60</sub>	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>61</sub>	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--	--
s <sub>70</sub>	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--	--
s <sub>71</sub>	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--	--
s <sub>80</sub>	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--	--
s <sub>81</sub>	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--	--
s <sub>90</sub>	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--	--
s <sub>91</sub>	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--	--
s <sub>100</sub>	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--	--
s <sub>101</sub>	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--	--
s <sub>110</sub>	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--	--
s <sub>111</sub>	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--	--
s <sub>120</sub>	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--	--
s <sub>121</sub>	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--	--
s <sub>130</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--	--
s <sub>131</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--	--
s <sub>140</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--	--
s <sub>141</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--	--
s <sub>150</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--	--
s <sub>151</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--	--
s <sub>160</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--	--
s <sub>161</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--	--
s <sub>170</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0	--
s <sub>171</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1	--
s <sub>180</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1	0
s <sub>181</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0	1
s <sub>190</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	1
s <sub>191</sub>	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	0

figure 5.2: Consolidated Satoku Matrix for Reduced CNF Formula Without At-Least-1 Clauses











Removing the mapped variables finally leaves an exact copy of the original message (see figure 5.8).

P	-----	-----	-----	-----	-----	-----
$s_{0_0}$	1 0 0 0 0 0	- 0 - - 0 -	- - 0 0 - -	- 0 0 0 - -	- 0 0 0 - -	- 0 - - - 0
$s_{0_1}$	0 1 0 0 0 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - 0 -	- 0 - - 0 -	- - 0 - 0 -
$s_{0_2}$	0 0 1 0 0 0	- 0 0 0 0 -	- 0 0 0 0 -	- 0 0 0 - -	- 0 0 0 - -	- - - 0 - -
$s_{0_3}$	0 0 0 1 0 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - - -	- 0 - - - -	- - - 0 - -
$s_{0_4}$	0 0 0 0 1 0	- 0 - - 0 -	- 0 - - 0 -	- 0 - - - -	- 0 - - - -	- - - 0 - -
$s_{0_5}$	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -	- - - - - -	- - - - - -
$s_{1_0}$	- - - - - -	1 0 0 0 0 0	- - 0 0 - -	- - 0 0 - -	- 0 0 0 - -	1 0 0 0 0 0
$s_{1_1}$	0 0 0 0 0 1	0 1 0 0 0 0	- 0 - - - -	- 0 - - 0 -	- 0 - - - -	- - - 0 - -
$s_{1_2}$	- - 0 - - -	0 0 1 0 0 0	- 0 0 - - -	- 0 0 0 0 -	- 0 0 0 0 -	- - - 0 - -
$s_{1_3}$	- - 0 - - -	0 0 0 1 0 0	- 0 - - - -	- 0 - - 0 -	- - - - 0 -	- - - 0 - -
$s_{1_4}$	0 0 0 0 0 1	0 0 0 0 1 0	- - 0 0 - -	- 0 - - 0 -	- 0 0 0 0 -	- - - 0 - -
$s_{1_5}$	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -	- - - - - -
$s_{2_0}$	- - - - - -	- - - - - -	1 0 0 0 0 0	- - 0 0 - -	- 0 0 0 - -	- - - 0 - -
$s_{2_1}$	- 0 0 0 0 -	- 0 0 0 - -	0 1 0 0 0 0	- 0 - - - -	- 0 - - 0 -	- - - 0 - -
$s_{2_2}$	0 - 0 - - -	0 - 0 - 0 -	0 0 1 0 0 0	- 0 0 - - -	- 0 0 0 - -	- - - 0 - -
$s_{2_3}$	0 - 0 - - -	0 - - - 0 -	0 0 0 1 0 0	- 0 - - - -	- 0 - 0 - -	- - - - - -
$s_{2_4}$	- 0 0 0 0 -	- - - - - -	0 0 0 0 1 0	- - 0 0 - -	- 0 - - 0 -	- - - 0 - -
$s_{2_5}$	- - - - - -	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -	- - - - - -
$s_{3_0}$	- - - - - -	- - - - - -	- - - - - -	1 0 0 0 0 0	- - - - - -	- - - - - -
$s_{3_1}$	0 0 0 0 0 1	- 0 0 0 0 -	- 0 0 0 - -	0 1 0 0 0 0	- 0 - - 0 -	- - 0 0 - -
$s_{3_2}$	0 - 0 - - -	0 - 0 - - -	0 - 0 - 0 -	0 0 1 0 0 0	- - - - - -	- 0 - - 0 -
$s_{3_3}$	0 - 0 - - -	0 - 0 - - -	0 - - - 0 -	0 0 0 1 0 0	0 - - - - 0	- 0 - - 0 -
$s_{3_4}$	- 0 - - - -	- 0 0 0 0 -	- - - - - -	0 0 0 0 1 0	- 0 0 0 0 -	- - 0 0 - -
$s_{3_5}$	- - - - - -	- - - - - -	- - - - - -	0 0 0 0 0 1	- - - - - -	- - - - - -
$s_{4_0}$	- - - - - -	- - - - - -	- - - - - -	- - - 0 - -	1 0 0 0 0 0	- - - - - -
$s_{4_1}$	0 0 0 0 0 1	0 0 0 - 0 -	0 0 0 0 0 1	- 0 - - 0 -	0 1 0 0 0 0	- - 0 0 - -
$s_{4_2}$	0 - 0 - - -	0 - 0 - 0 -	0 - 0 - - -	- - - - 0 -	0 0 1 0 0 0	- 0 - - 0 -
$s_{4_3}$	0 - 0 - - -	0 - 0 - 0 -	0 - 0 0 - -	- - - - 0 -	0 0 0 1 0 0	- 0 - - 0 -
$s_{4_4}$	- 0 - - - -	0 - 0 0 0 -	- 0 - - 0 -	- 0 - - 0 -	0 0 0 0 1 0	- - 0 0 - -
$s_{4_5}$	- - - - - -	- - - - - -	- - - - - -	- - - 0 - -	0 0 0 0 0 1	- - - - - -
$s_{5_0}$	- - - - - -	- - - - - -	- - - - - -	- - - - - -	- - - - - -	1 0 0 0 0 0
$s_{5_1}$	0 - - - - -	0 - - - - -	- - - - - -	- - 0 0 - -	- - 0 0 - -	0 1 0 0 0 0
$s_{5_2}$	- 0 - - - -	0 - - - - -	- - - - - -	- 0 - - 0 -	- 0 - - 0 -	0 0 1 0 0 0
$s_{5_3}$	- - 0 0 0 -	0 0 0 0 0 1	0 0 0 - 0 -	- 0 - - 0 -	- 0 - - 0 -	0 0 0 1 0 0
$s_{5_4}$	- 0 - - - -	0 - - - - -	- - - - - -	- - 0 0 - -	- - 0 0 - -	0 0 0 0 1 0
$s_{5_5}$	0 - - - - -	0 - - - - -	- - - - - -	- - - - - -	- - - - - -	0 0 0 0 0 1

figure 5.8: Removing Mapped Variables Leaves Exact Copy of Original Message



